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Forensic Analysis of IoT File Systems for Linux-Compatible Platforms

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Abstract: Due to recent developments in IT technology, various IoT devices have been developed for use in various environments, such as card smart TVs, and smartphones Communication between IoT devices has become possible. Various IoT devices are found in homes and in daily life, and IoT technologies are being combined with vehicles, power, and wearables, amongst others. Although the usage of IoT devices has increased, the level of security technology applied to IoT devices is still insufficient. There is sensitive information stored inside IoT devices, such as personal information and usage history, so if security accidents happen, such as data leakage, it can be very damaging for users. Since research on data storage and acquisition in IoT devices is very important, in this paper we conducted a security analysis, from a forensic perspective, on IoT platform file systems used in various environments. The analysis was conducted on two mechanical platforms: Tizen (VDFS) and Linux (JFFS2 and UBIFS). Through file system metadata analysis, file system type, size, list of files and folders, deleted file information were obtained so that we could analyze file system structure with the obtained information. We also used the obtained information to check the recoverability of deleted data to investigate the recovery plan. In this study, we explain the characteristics of platforms used in various environments, and the characteristics of data stored in each platform. By analyzing the security issues of data stored during platform communications, we aimed to help in solving the problems affecting devices. In addition, we explain the analysis method for file system forensics so that it can be referred to in other platform forensics.

Keywords: digital forensic; IoT devices; VDFS; JFFS2; UBIFS



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1. Introduction

As the Internet of Things (IoT) advances and the market size increases, increasingly advanced technologies are being applied to various devices [1]. As of 2022, IoT technologies are being used in various devices and contexts, including vehicles, artificial intelligence, wearables, speakers, and industrial products, and it is becoming increasingly difficult to find devices without these technologies [2,3]. IoT devices store a variety of data acquired through interactions with networks, devices, and people. Such stored data includes sensitive user information, including system logs and personal information [4].

However, the level of security applied to IoT devices remains insufficient and devices such as speakers and intercoms are vulnerable to various attacks. In one case, an individual's IP camera was hacked, and the resulting video was sold to an Internet site [5]. IoT devices have become ubiquitous, and significant harm can occur through attacks. To obtain proper evidence, forensic and security technologies need to be developed [6–8].

As the fields in which IoT devices and technologies are applied continue to diversify, the numbers of IoT platforms are increasing. In addition, because each platform has a different data storage method and structure, knowledge of various platforms is required. Although research has been conducted on certain platforms, there are still platforms for which no studies have yet been conducted [9–12].

In addition, the types of IoT platforms and security technologies of the platforms are developing daily [13]. Due to the development of technologies, it is difficult to obtain data from IoT devices because of improved security and encryption. Therefore, research is needed on data acquisition and analysis methods in various environments to investigate and analyze devices.

Platforms are mounted on various IoT devices, so such devices can be targets for attack. IoT devices can store important information, such as user activity and usage records [14]. It is, therefore, necessary to research platforms to improve the security of IoT devices and the efficiency of evidence collection. With this in mind, we analyzed the metadata of file systems on various platforms for improved security and evidence collection. In addition, we propose a method to extract data through analyzed metadata and a method to recover deleted data. Finally, we investigate the types of artifacts that can be obtained for each file system by applying the proposed forensics method. The contributions of this study are as follows:

- Through a metadata analysis of the file systems mounted on IoT devices, information, such as the system type and size, list of files and folders, and deleted files, was obtained.
- The possibility of extracting files stored in a file system was analyzed experimentally. The metadata on devices using the Vertically Deliverable File System and Journaling File system version 2 (JFFS2) were analyzed to confirm the possibility of data extraction through file carving.
- The recoverability of deleted files in a file system was analyzed experimentally. Data
 on deleted files were confirmed using the snapshot and journaling functions applied
 by the file system.
- Analysis and file recovery experiments were conducted on real devices and it was confirmed that the research technique could be applied to other real IoT devices.

By analyzing the data storage structure and security issues of the platform, the security of IoT devices to be developed in the future can be improved. In addition, we present an approach for analyzing the platform mounted on IoT devices using file system forensics. The proposed approach allows for referencing when analyzing IoT device platforms that are newly developed or difficult to recognize.

The remainder of this paper is structured as follows. Section 2 describes existing studies related to file system forensics on various platforms. Section 3 details the characteristics of each platform and applicable devices, and Section 4 introduces the metadata analysis experiments conducted using a file system. Section 5 describes the extraction and recovery of files based on case studies. Section 6 discusses the results of such a study, and, finally, Section 7 provides some concluding remarks.

2. Related Work

IoT file system forensics requires both analysis of characteristics of IoT devices and knowledge of file systems. Since IoT has different characteristics for each device, a forensic framework that can combine various devices is needed.

In 2019, Jianwei Hou et al. confirmed the importance of forensics for IoT devices and investigated papers that conducted forensics for various IoT devices. Classifying the papers into forensic techniques, the study made it easier to refer to the necessary parts in future research [15].

In addition, other studies investigated various devices to which IoT technology can be applied and analyzed the problems facing IoT forensics. Through this, it was possible to identify devices to which IoT forensics can apply and to gauge the direction in which forensics technology should go [16,17].

In 2019, Arduino's research was conducted to simulate attacks that may occur on IoT devices, and to detect and identify attacks. Since this study simulated attacks that could easily occur in a network environment, the results of the study could help in the event of a real problem [18].

In 2022, the environment of IoT devices was formed through Raspberry Pie and forensics was conducted. This study, conducted on IoT, attempted to acquire data stored and communicated through connection with the network [19]. As such, various forensic studies were conducted on IoT devices, and could be applied to various devices.

A file system refers to a method of storing, maintaining, and managing data on a device. There are various types of file systems, and the method of managing data and the structure of storing data are all different for each file system. It is difficult to confirm the existence of a file system as a general user who uses only IoT devices. However, the file system contains a lot of information. As an example, there is information on deleted data from the stored data, which cannot be generally confirmed. To confirm this, metadata can be obtained through file system analysis, and various data can be accessed through the obtained information [20–22].

Since file systems are mounted on a wide variety of devices, there are various types, from file systems installed on basic computers and smartphones to file systems installed on small IoT devices [23,24]. As new types of devices are developed, the number of file systems suitable for devices is expected to continue to increase.

Research on file systems is being conducted, and file system forensics research is also being actively conducted in various fields, such as computers, smartphones, and IoT.

In 2019, a study was conducted to extract the files of a Samsung smart TV [25]. This study obtained and analyzed the data and extracted files through the chip-off technologies used by Samsung smart TVs and flash memory. However, the study only analyzed the file extraction and artifacts, and no research was conducted on Samsung's own file system.

In 2019, a forensic experiment was conducted on Amazon Echo [26]. The study applied a basic forensics model, and an analysis was conducted by dumping the firmware through UART communication. Based on a firmware analysis, various types of information, including the data type and MAC address stored in Amazon Echo, was obtained. However, this study only introduced an approach for a firmware analysis.

In 2019 and 2022, data analysis and extraction of data from various wearable devices were conducted [27,28]. This study considered the platforms of various wearable devices and organized the data stored on the platform, as well as the data acquisition. Although this study intensively analyzed the available evidence and file acquisition on various wearable device platforms, the file system layer was not considered.

In 2022, a study on smart home forensics using IoT was conducted [29], and the results indicated that data stored on a smart phone can be confirmed during communication with smart home IoT devices. However, only data stored on a smartphone, and not a smart-home IoT device, were analyzed.

In 2022, Android smartphones were analyzed in terms of the recovery and exportation of deleted data and a file recovery of the ext file systems [30]. In this study, forensics was conducted through metadata analysis of the file system used in Android smartphones.

In 2022, the application of IoT technology to vehicles, prompted forensic research on Android auto and apple CarPlay [31]. This study focused on the connection between devices and clouds, IVI devices and mobile devices, that is, network communication. However, it is significant in that it attempted to analyze new devices because analysis of IVI devices and mobile devices is necessary to obtain communication data stored in the device. In addition, research on data acquisition was conducted on Xiaomi Mi Smart Sensor, but not much on the file system stage [32]. It was a study on data recovery through the analysis of Google home and the analysis of SquashFS, a file system, and had a similar purpose to this paper. With the existence of forensic papers targeting various devices, it is possible to confirm the necessity of analysis at the file system stage of the IoT platform [33].

In addition to research on devices and platforms, framework research for forensic techniques has also been continuously studied [34–36]. However, since IoT devices are equipped with various platforms, further research on unresearched platforms is needed.

As such, research on IoT platforms is steadily progressing, but it is still insufficient. Some platforms have been analyzed at the file system stage, but there are still many

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platforms that require further research. Related studies have focused on acquiring data from peripheral devices or analyzing files extracted from the device itself. However, it may be difficult to check deleted files by this method, and information on the acquired file may be limited.

To solve this problem, there is a need for a file system-based forensics technique that can obtain various data including deleted data. Although some file system forensics techniques have been studied, continuous research is required, because the characteristics of each platform or file system are different [37]. Therefore, in this study, a file system analysis was conducted to overcome these shortcomings.

3. IoT Platforms

In this section, we describe the characteristics of the platforms and file systems used by IoT devices. Various types of M2M equipment are currently being used, and so various platforms are also applied. In addition, the operating and file systems may differ, even for the same platform. Platforms applied to such devices include Fuchsia, Contiki, Tizen, and Linux.

Fuchsia is a platform developed by Google [38]. Fuchsia OS is currently experimentally installed in Google Nest Hub (1st Gen) and uses BlobFS and MinFS as the file systems. MinFS is a file system built for the Zircon kernel applied by Fuchsia. Fuchsia is expected to replace Android in the future.

Platforms are also used for low-power and low-capacity IoT devices. The Contiki platform is applied in low-power devices and uses Contiki OS and the coffee file system [39]. Contiki is run in memory, and because, it is mainly applied to sensors, lighting, and monitoring devices in smart cities, it is likely to be used in the future.

With the application of various platforms, to achieve improvements in forensic analysis, the necessity of conducting a platform analysis has been confirmed. Therefore, in this study, we conducted an experimental forensic analysis on two types of platforms: Tizen and Linux. The basic structures of the file systems used in Tizen and Linux was verified for this experiment.

3.1. Tizen (VDFS)

Tizen is a Linux-based platform developed by Linux, Samsung Electronics, and Intel, in 2012. Devices using the Tizen platform are primarily equipped with Tizen OS. Although various file systems can be used for Tizen OS, Samsung's own file system, VDFS, was applied during the experiment. VDFS is a Linux file system optimized for usage on eMMC flash devices. VDFS is a file system produced by Samsung that provides code as an open source by Samsung, but not much information is known about the file system. Through the analysis of the code provided by open source, it was found that VDFS used the btree structure when storing files, and deleted files were stored in the form of snapshots, rather than journal areas. Therefore, we could analyze the btree structure to extract stored files, and deleted files could also be extracted or recovered if snapshot data was not erased. Although VDFS is only used by Samsung, we conducted research on this file system owing to the high penetration rate of Samsung TVs. Figure 1 shows the system–layer structure of Tizen VDFS, which is configured in a manner similar to that of a basic Linux file system. VDFS operates on block devices, stores data in a block-based manner, and uses a btree structure for storage.

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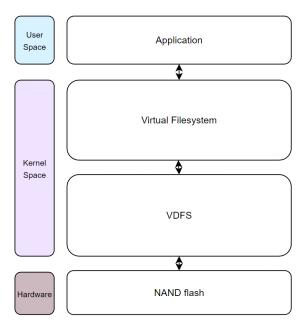


Figure 1. VDFS architecture on Tizen.

3.2. Linux (JFFS2 and UBIFS)

Linux is an open-source platform developed in 1991 and is mounted on various devices and PCs. Although the Linux OS used on the Linux platform can apply a variety of file systems, JFFS2 and UBIFS were used in the experiments. JFFS2 and UBIFS are Linux-based file systems used in flash memory devices. When storing files, JFFS2 can store files using three types of compression algorithms, zlib, rubin, and rtime, to manage data capacity more efficiently. Since JFFS2 is a journaling file system, files can be backed up and stored, and it is possible to recover the file even if the file is deleted later. UBIFS was developed based on improvements to JFFS2, and unlike JFFS2, it does not scan the entire media, and can, thus, be used for large capacity NAND flash memory. In addition, since JFFS2 stores file system indexes in memory, while UBIFS stores indexes in flash, there is a difference in performance and speed. When storing files, UBIFS encrypts file data through fcrypt, making it impossible to extract general files. When storing files, UBIFS encrypts file data through fscrypt, making it impossible to extract files in a conventional way. Therefore, with UBIFS using the journaling system, deleted files also have data, but it is difficult to recover deleted files due to fscrypt encryption. These file systems are mainly used in embedded systems, and although many years have passed since their initial development, they are still applied in certain models. Therefore, we conducted a metadata analysis and file extraction for the file system. Figure 2 shows the system-layer structures of JFFS2 and UBIFS. JFFS2 operates on memory technology device (MTD) devices, but UBIFS operates on UBI volumes and only on raw flash memory. For UBIFS, the UBI volume is used to access the MTD devices. In addition, UBIFS store data using a logical erasable block (LEB). Electronics **2022**, 11, 3219 6 of 19

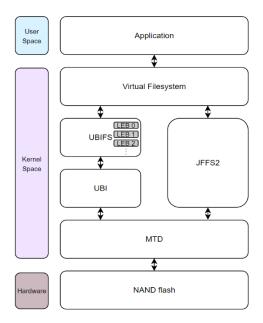


Figure 2. JFFS2 and UBIFS architecture on Linux.

4. Experiments and Analysis

In this section, experiments conducted to acquire data from a device are described, along with an analysis on the platform and file system used.

4.1. Experiments on Tizen Using VDFS

For a VDFS analysis of the Tizen platform, a study was conducted on a smart TV (model KU43T5300AFXKR). An analysis of the PCB of the model, as shown in Figure 3, revealed that a THGBMNG5D1LBAIT NAND flash chip with 4 GB of memory and 153 BGA was used. The chip-off technique was applied to obtain data for the corresponding NAND flash. Raw data was obtained through a DS3000-USB3.0-emmc153 reader, a product of Allsocket in Dongguan, China. Based on an analysis of the raw data obtained through the FTKimager analysis tool (v4.3.0.18), 26 partitions were recognized, as shown in Figure 4. Among the 26 partitions, 5 were confirmed to be VDFS partitions with VDFS 2007 magic numbers. A file system analysis was conducted only on the VDFS partitions.



Figure 3. Nand flash chip (KU43T5300AFXKR) present on pcb of smart TV.

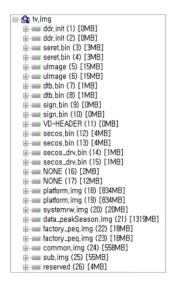


Figure 4. Smart TV partition.

For the analysis of the VDFS file system, we analyzed the code of the VDFS released by Samsung and proceeded. VDFS, a file system manufactured by Samsung, has a magic number labeled VDFS2007 at the front address of the partition, indicating that the partition was a VDFS. The superblock could be checked at the 0x400 location of the VDFS partition. The structure of the superblock is shown in Figure 5, and basic information on the partition, such as the version, creation time, and volume name, could be confirmed. Behind the superblock was an extended superblock with additional information at the 0x600 location. The structure of the extended superblock is as shown in Figure 6, and information on the stored data, such as the number of files, folders, and meta tree offset in the volume, could be found.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	ос	0D	OE	OF
00		signa	ature			vers	sion				max	imum l	olock c	ount		
10			creation timestamp													
20					cre	ation t	imesta	mp						volum	e uuid	
30						volum	e uuid							volume	name	
40						volume	e name							mkfs v	ersion	
50								mkfs v	ersion				•			
Α0					log blk size											

Figure 5. VDFS superblock structure in form of hex value.

An analysis of a superblock allowed us to obtain the offset of the meta tree, and when we moved to the acquired address, the xattr tree, catalog tree, inode bitmap, and fsm bitmap were located. The tree haf a magic number of 0x644E at the beginning of the block, and bitmap had magic numbers of 'inob' and 'fsmb'. In addition, the xattr tree had 'XAre' as its magic number. The xattr tree stored records such as the extension properties and access control, rather than regular files, and the catalog tree stored records of the files and folders. We could check the file and folder names in the catalog tree. Figure 7 shows a record of the airplay-service.md file in the catalog tree.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	ос	0D	OE	OF
00				file c	ount							folder	count			
10				volum	e start							volum	e end			
20	1	mount	counte	r		sync c	ounter		u	mount	count	er		eneration mber g		
30								debu	g area							
40	bt	ree ext block	ents to count	tal		pade	ding					tab	oles			
50				tab	les						meta	a btree	start o	offset		
60			m	eta btre	ee leng	th										

Figure 6. VDFS extended superblock structure in form of hex value.

```
1CE3A000 4E 64 CA 74 A8 01 00 00 06 00 00 00 A0 11 60 00
                                                1CE3A010 DF 01 00 00 DE 01 00 00 E2 01 00 00 02 00 00 00
                                                ß...Þ...â.....
1CE3A020 00 00 00 00 30 00 30 00 22 41 00 00 00 00 00
                                                ....0.0."A....
1CE3A030 20 41 00 00 00 00 00 05 12 61 69 72 70 6C 61
                                                 A....airpla
1CE3A040 79 2D 73 65 72 76 69 63 65 2E 6D 6F 00 00 00 00 y-service.mc....
1CE3A050
       00 00 00 00 28 00 78 00 24 41 00 00 00 00 00 00
                                                 ....(.x.$A....
1CE3A060 26 41 00 00 00 00 00 01 0B 4C 43 5F 4D 45 53
                                                &A....LC_MES
1CE3A070 53 41 47 45 53 00 00 00 00 00 00 04 00 00 00
                                                SAGES.....
......
1CE3A090 FF FF FF FF FF FF FF ED 41 15 13 C9 00 00 00
                                                ÿÿÿÿÿÿÿÿiA..É...
```

Figure 7. The airplay-service.md file records in a catalog tree with magic number 0x644E.

The catalog tree contained all the information on the files and folders. Figure 8 shows the structure of a file in the catalog tree. The type of data could be checked through the variable value 'type'. A value of 'type' 0x01 indicated a folder; 0x02 was a file; and 0x05 was a link record. Each dataset consisted of a key + record, which contained basic information, such as the name, type, and object ID. The record contained the size, creation/modification/access time, and offset address information of the file where the data were stored. However, if the type was a folder, the number of files present in the folder was written instead of the file size. If the type was a link record, only a key value was present.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	0C	0D	0E	OF
00	key	len	recor	d len				obje	ct id					pare	nt id	
10		pare	nt id		type	File name len			file r	name				-		key len end
20			-				•	file size	e (byte)							
30									-							
40					-							creatio	n time			
50		creatio	n time						m	odifica	tion tir	ne				
60						access	time							file size	(byte)	
70		file size	e (byte)			nun	nber of	f all blo	cks				data	offset	
80		data	offset				nur	mber o	f block	s_1				start	block	
90		start	block					data o	ffset_2				nui	mber o	f block	s_2
Α0	nui	mber o	f block	:s_2				start b	lock_2							

Figure 8. Catalog tree data structure in form of hex value.

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VDFS used snapshots to store the file system formation. Snapshots store data in files and folder-like images. Figure 9 shows the 'browser-data.db' file, and we confirmed that the offset addresses of the stored data differed because the contents changed as the 'browser-data.db' file was modified, and the previous files were saved as snapshots. If we looked at the offset corresponding to each file, data in the 'browser-data.db' file were present at that time. In this manner, we could extract or recover files using the snapshot function. Therefore, experiments related to file extraction and recovery were conducted to analyze this possibility, and the effects were verified through a case study.

```
1D30A450
         00 00 00 00
                   30
                         68 01
                              03
                                 33 00
                                         00 00
1D30A460
         04 33 00 00 00 00 00 00 02 10 2E 62 72 6F 77 73
                                                     er-data.db.....
1D30A470
         65 72 2D 64 61 74 61 2E 64 62 00 00 00 00 00 00
                                                     .....p.....p
1D30A480
         00 00 00 00 04 00 00 00 00
                                 70 01 00 00 00 00 00
1D30A490
        1D30A4A0
        A4 81 02 00 C9 00 00 00 00 00 00
                                      00 68 C2 6F 60
                                                     ¤...É.....hÂo
1D30A4B0
        00 00 00 00 AD 73 2C 16 57 C2 6F 60 00 00 00 00
                                                     ....s,.WÂo`...
1D30A4C0
        B1 06 B9 05 CD 1F 79 60 00 00 00 00 34 D3 F2 01
                                                     ±.1.Í.y`....4Óò.
1D30A4D0
                   00
                         00
                              17 00 00 00 00 00 00 00
                                                     .p.........
                         00 00 16 00 00 00 00 00 00 00
1D30A4E0
              00 00 00 00
1D999640
         00 00 00 00 30 00 68 01 1A 1F 00 00 00 00 00 00
                                                      ...0.h......
1D999650
        1B 1F 00 00 00 00 00 00 02 10 2E 62 72 6F 77 73
10999660
         65 72 2D 64 61 74 61 2E 64 62 00 00 00 00 00 00
                                                     er-data.db....
                                                     ....p.....
1D999670
        00 00 00 00 05 00 00 00 00 70 01 00 00 00 00
1D999680
        •••••
10999690
        A4 81 00 00 89 13 00 00 1B 04 00 00 CD 1F 79 60
                                                     ¤.....í.y
1D9996A0
        00 00 00 00 34 27 DD 2C CD 1F 79 60 00 00 00 00
                                                      ....4'Ý,Í.y`...
1D9996B0
        34 27 DD 2C CD 1F 79 60 00 00 00 00 34 30 1A 2D
                                                     4'Ý,Í.y`....40.-
1D9996C0
                 00 00 00 00 00 17 00 00 00 00 00 00 00
1D9996D0
                         00 00 17 00 00 00 00 00 00 00
```

Figure 9. Two different offsets (blue) of 'browser-data.db' file (red).

4.2. Experiments on Linux Using UBIFS/JFFS2

In this study, we investigated the data central units (DCUs) of LOENK, a power device product. As shown in Figure 10, 256 MB NAND flash with a K9F2G08UOC chip was used. To obtain data from NAND flash, serial communication was executed through the RS-232 port of the PCB. As a result of accessing and checking the root account through serial communication, three MTDs were identified. After creating an SSH server for the extraction of the MTD, we dumped the MTD using a scp command to transfer files through SSH on the device. As a result of analyzing the MTD, it was confirmed that a kernel, JFFS2, and UBIFS were installed. The file systems were also analyzed.

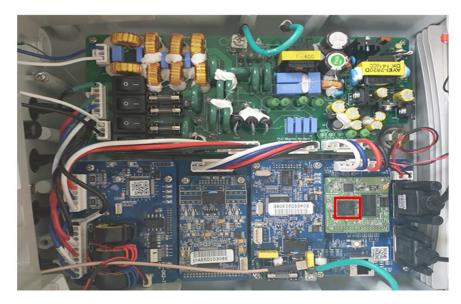


Figure 10. PCB of DCU.

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For the analysis of JFFS2 and UBIFS, an analysis of the code disclosed as an open source was conducted. JFFS2 had a magic number labeled 0x1985 at the front of the node, and the unused area was filled with 0xFF values. JFFS2 consisted of several node types. Among the nodes, Figure 11 shows the structure of the inode, which had a node type of 0xE002. In addition, the inode contained various information on the file, including file size, modification, access time, and data. JFFS2 compressed and stored data for efficient data management, and only zlib, rubin, rtime, and LZO were used for the compression algorithm.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	ос	0D	0E	OF
00	0x1	985	node	type		node	e_len			head	er_crc		i	inode r	number	
10		vers	sion			file r	node		ui	id	g	id		isi	ze	
20		ati	me			mti	me			cti	me			Off	set	
30	C	ompres	ssed siz	ze	de	compre	essed s	ize	comp	usr comp	fla	gs		data	_crc	
40		node	e_crc							da	ita					

Figure 11. JFFS2 inode structure in form of hex value.

In JFFS2, a dirent node indicated the inode. Figure 12 shows the structure of the dirent node, and the corresponding node showed the file name and modified time stored in the inode. Based on the inode number of the corresponding node, an inode had the same inode number value, and the location where data existed could, thus, be determined.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	0C	0D	0E	OF
00	0x1	985	node	type		node	e_len			head	er_crc			parent	inode	
10		ver	sion			inode r	numbei	r		mct	ime		name size	type	unu	sed
20		node	e_crc			nam	e_crc					na	me			

Figure 12. JFFS2 dirent node in form of hex value.

Next, UBIFS was analyzed. UBIFS operates through the UBI volume, sp the UBIFS region had to be extracted from this volume if data from the MTD were to be obtained. In the UBI, an EC header containing information on the volume for each block existed, and in the EC header, basic information, such as the number of errors and the start offset of the data, was stored. In addition, we could move to the vid header through the vid header offset of the EC header. The vid header stored information on volumes that did not exist in the EC header, such as the LEB size and number of LEBs used in the volume.

UBIFS image extraction was possible using Github ubi reader (https://github.com/jrspruitt/ubi_reader, accessed on 5 May 2022). In this study, UBIFS images were extracted using ubi reader v0.8.0 in a Linux environment. UBIFS was in the form of a node. Figure 13 shows the structure of a header, which is a common form of all nodes, and all nodes in UBIFS had 0x06101831 as their magic number placed in front.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	ос	0D	OE	OF
00	-	magic r	numbe	r		CI	rc					sqn	um			
10		le	n		node type	group type	pado	ding				-	-			

Figure 13. UBI common header in form of hex value.

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The first node in UBIFS was a superblock. In a superblock node, basic information, such as the LEB size and LEB number of UBIFS, could be found; in addition, the LEB size became important when analyzing the metadata.

The second LEB had a master node. Figure 14 shows the structure of the master node, which was used in a similar manner as the superblock node. Information on the volume might also be stored, such as the size of the available area and the number of LEBs used. The master node also contained information regarding the location of the index node of the root and we could use this information to find the index node.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	ос	0D	OE	OF
00					-							highes	t_inum			
10				cmt	_no					fla	gs			log_l	num	
20		root_	lnum			root	_offs			root	_len			gc_lı	num	
30		ihead	_lnum			ihead	l_offs					Index	_size			
40				total	_free							total_	_dirty			
50				total_	used							total_	dead			
60				total	_dark					lpt_l	num			lpt_	offs	
70		nhead	l_Inum			nhead	d_offs			ltab_	lnum			ltab_	offs	
80		lsave_	_lnum			Isave	_offs			lscan	_Inum			empty	/_lebs	
90		ldx_	lebs			leb_	_cnt					pade	ding			

Figure 14. UBIFS master node in form of hex value.

Index nodes had information on the LEB numbers and node offsets, and when moving along the offset pointed to by a node, nodes with various information regarding the files, including the inode, data nodes, and directory entry nodes, could be accessed. Directory nodes had inode numbers for the files and folders, and inode and data nodes had data on the files. The structure of a data node, which stored data on the file, is shown in Figure 15.

	00	01	02	03	04	05	06	07	08	09	0A	ОВ	ос	0D	OE	OF
00					-							ke	еу			
10				ke	Э у					si	ze		con ty	npr_ pe		npr_ ze
20								da	ita							

Figure 15. UBIFS data node in form of hex value.

5. Case Studies

In this section, we analyze a file system based on a case study. The case study proceeded with the file extraction and recovery of deleted files using the file system. The file extraction proceeded with file carving by searching for file data through a metadata analysis. The recovery of deleted files proceeded by comparing such files with areas where data existed.

5.1. File Extraction

5.1.1. VDFS

To extract data from VDFS, a catalog tree had to be found. To access the catalog tree, we had to find the meta tree start offset in the extended superblock. The xattr tree, catalog tree, inode bitmap, and fsm bitmap were located in this offset. We could find a catalog tree to check the file record. Figure 16 shows the record for the 8669.png file, as well as information on the file name, file modification/generation time, and offset where the file data were stored. Information was also available on the total number of blocks where the data were stored, as well as the number of blocks stored in each offset.

									File	nam	e		Nu	ımbe	er of	all b	locks 🔲 Data offset
									File	size			Nu	mbe	er of	bloc	ks_1
1E7AE5A0	00	00	00	00	28	00	60	01	D2	25	00	00	00	00	00	00	(.`.Ò%
1E7AE5B0	E0	2F	00	00	00	00	00	00	02	08	38	36	36	39	2E	70	à/8669.p
1E7AE5C0	6E	67	00	00	00	00	00	00	00	00	00	00	0C	00	00	00	ng
1E7AE5D0	6F	39	01	00	00	00	00	00	01	00	00	00	00	00	00	00	09
1E7AE5E0	FF	FF	FF	FF	FF	FF	FF	FF	A4	81	00	00	89	13	00	00	ΫΫΫΫΫΫΫΫ×··.‱···
1E7AE5F0	64	00	00	00	D9	D2	DA	60	00	00	00	00	EC	B5	22	18	dὑο΄΄ìμ".
1E7AE600	D9	D2	DA	60	00	00	00	00	EC	B5	22	18	D9	D2	DA	60	ÙÒÚ`ìμ".ÙÒÚ`
1E7AE610	00	00	00	00	EC	82	70	02	6F	39	01	00	00	00	00	00	ì,p.o9
1E7AE620	14	00	00	00	00	00	00	00	DD	В3	00	00	00	00	00	00	ݳ
1E7AE630	14	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
1E7AE640	00	00	00	00	00	0.0	00	00	00	00	00	00	00	00	00	00	

Figure 16. File record representing metadata where the 8669.png file is stored.

Therefore, the offset in which the data were stored might be known by multiplying the value (0xB3DD) representing the data offset based on the block size. When moving to the offset, it could be confirmed that data on the 8669.png file existed, as shown in Figure 17. Thereafter, a file could be extracted by performing data carving on an area as large as the data size.

```
FF D8 FF E0 00 10 4A 46 49 46 00 01 01 00 00 01
0B3DD000
                                                          ÿØÿà..JFIF.....
OB3DD010 00 01 00 00 FF DB 00 43 00 03 02 02 03 02 02 03
                                                         ....ÿÛ.C.....
OB3DD020 03 03 03 04 03 03 04 05 08 05 05 04 04 05 0A 07
                                                          . . . . . . . . . . . . . . . .
OB3DD030 07 06 08 0C 0A 0C 0C 0B 0A 0B 0B 0D 0E 12 10 0D
                                                          . . . . . . . . . . . . . . . .
         OE 11 OE 0B OB 10 16 10 11 13 14 15 15 15 OC OF
0B3DD040
                                                          . . . . . . . . . . . . . . . . .
0B3DD050
         17 18 16 14 18 12 14 15 14 FF DB 00 43 01 03 04
                                                          ....ÿÛ.C...
         04 05 04 05 09 05 05 09 14 0D 0B 0D 14 14 14 14
0B3DD060
                                                          . . . . . . . . . . . . . . . .
0B3DD070
         . . . . . . . . . . . . . . . .
0B3DD080
         . . . . . . . . . . . . . . . .
```

Figure 17. 8669.png file data.

5.1.2. JFFS2

To extract a file from JFFS2, it was necessary to determine the inode of the file. However, because file name data did not exist in the inode, we needed to find the dirent node to check the file name. A file might be extracted by finding an inode with the same inode number as the rent node. However, because JFFS2 compressed files when they were stored, file extraction required decompression. Figure 18 shows the dirent node and inode of test3.txt. The file name was located in the dirent node, and data could be found in the inode.

] C	omp	resse	d size	е		No	de typ	e [File name
] D	econ	npres	sed s	ize		Co	mpres	sion type	data
01A95040	6A	4D	6A	DC	85	19	01	E0	31	00	00	00	1D	D9	82	42	jMjÜ	àlÜ,B
01A95050	01	00	00	00	07	00	00	00	06	00	00	00	63	16	C2	62		c.Âb
01A95060	09	08	00	00	DF	ED	00	El	6B	25	6D	3D	74	65	73	74	Bí.	ák%m=test
01A95070	33	2E	74	78	74	FF	FF	FF	85	19	02	E0	6A	00	00	00	3.txtÿÿ	ÿàj
01A95070	33	2E	74	78	74	FF	FF	FF	85	19	02	ΕO	6A	00	00	00	3.txtÿÿ	ÿàj
01A95080	10	23	1A	D8	06	00	00	00	02	00	00	00	A4	81	00	00	.#.Ø	¤
01A95090	00	00	00	00	57	01	00	00	63	16	C2	62	63	16	C2	62	W	.c.Âbc.Âb
01A950A0	63	16	C2	62	00	00	00	00	26	00	00	00	57	01	00	00	c.Âb	. & W
01A950B0	06	00	00	00	9C	FA	27	7A	6A	FA	AF	57	78	5E	2A	C9	Ϝ'	zjú Wx^*É
01A950C0	C8	2C	56	00	A2	92	D4	E2	12	63	BD	92	8A	12	AE	92	È, V. c' Ĉ	â.c½'Š.⊗'
01A950D0	51	11	F2	42	03	20	80	90	74	71	01	04	18	00	CC	6A	Q.òB. €	.tqÌj
01A950E0	79	18	FF	FF	85	19	02	E0	44	00	00	00	1D	FB	F7	98	у. ÿÿ	àDû÷~

Figure 18. Dirent node (top) and inode (bottom) of 'test3.txt' file.

Although data were present in the inode, they were compressed, and the compression type could be checked. The compression type was mainly 0x06, which indicated that the data had been compressed using zlib. In order to confirm the contents of the compressed data, decompression should be performed in the same manner as the compressed type. The data was compressed into zlib, and in the case of zlib, it could be decompressed through a simple python code. Figure 19 is a code for decompressing a file compressed with zlib. The data size of the compressed file could be found through the compressed size metadata, and the original data could be obtained through the code of Figure 19 after carving the data area. The size of the decompressed file could be compared with the decompressed size metadata to verify that the file had been decompressed properly.

```
import zlib
file = open('test3', 'rb')
data = file.read()
file.close()
d_data = zlib.decompress(data)
print(d_data)
b'this is test3.txt\nthis is test3.txt\nthis is test3.txt\nthis
```

Figure 19. Zlib Python code to decompress files and the result of decompression.

5.1.3. UBIFS

A data node had to be accessed to extract the UBIFS files. To do this, we needed to move to the master node, which was placed in the second LEB. The indexing node offset was obtained and we moved through an analysis of the master node. We could then follow the branch of the indexing node to observe the directory entry and inode of the files and folders in that directory. Figure 20 shows the directory entry and inode, which can be distinguished by the node type value.

By searching for the inode number of the file, the file data node could be found. In this case, the last value of the inode number of the data node was changed to 0x20. Figure 21 shows the directory entry and data node of the k2logd.sh file. By comparing the inode number values, we could see that it was the same file.

However, unlike JFFS2, UBIFS applied compression and encryption in the data area, which made it difficult to extract the files. Figure 22 shows the data of the dcu.tar file and the data node of this file stored in the UBIFS. It can be seen that these data differed from existing data owing to the encryption and compression of the UBIFS. Thus, in the case of a UBIFS file, decoding was required for extraction.

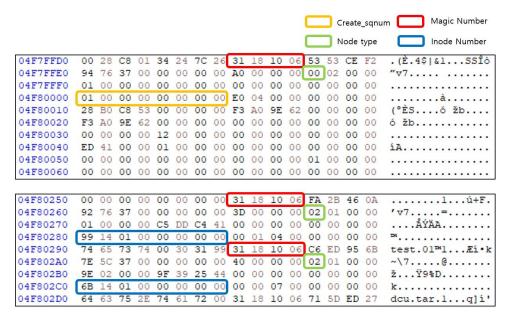


Figure 20. UBIFS directory of entry (top)/inode (bottom).

															ata		Magic Number
														r	lode	type	Inode Number
004FBE90	31	18	10	06	7A	Fl	8E	48	87	OF	00	00	00	00	00	00	1zñŽH‡
004FBEA0	42	00	00	00	02	00	00	00	EA	02	00	00	AC	9F	0B	4F	B꬟.O
004FBEB0	00	00	00	00	00	00	00	00	F4	02	00	00	00	00	00	00	ô
004FBEC0	00	00	09	00	00	00	00	00	6B	32	6C	6F	67	64	2E	73	k2logd.s
004FBED0	68	00	FF	FF	FF	FF	FF	FF	31	18	10	06	95	E8	15	69	h.ÿÿÿÿÿÿl•è.i
004FBD80	31	18	10	06	EB	42	7E	72	85	OF	00	00	00	00	00	00	1ëB~r
004FBD80 004FBD90	31 70	18	10 00	_	EB 01		7E 00		85 F4	-	00	00	00	00	00	00 20	1ëB~r pô
	31 70 00	18 00 00		_	and the latest designation of		1.000			-		00		00			
004FBD90	31 70 00 23	00	00	00	01	00	00	00	F4	-		00 00 00 77		00 00 00 69	00	20 00	pô
004FBD90 004FBDA0	70 00	00 00	00 00 2F	00	01 00	00	00	00 00 73	F4 40	02 00	00	00	00	00	00 00 6C	20 00	pô
004FBD90 004FBDA0 004FBDB0	70 00 23	00 00 21	00 00 2F	00 00 62	01 00 69	00 00 6E	00 00 2F	00 00 73	F4 40 68	02 00 20	00 00 0A	00 77	00 00 68	00 69	00 00 6C	20 00 65	pô @ #!/bin/sh .while

Figure 21. UBIFS directory of entry (top)/data node (bottom).

									C		Ma	agic I	Numb	oer			Node type Data
00000000	64	63	75	2F	00	00	00	00	00	00	00	00	00	00	00	00	dcu/
00000010	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000030	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000040	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000050	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000060	00	00	00	00	30	30	30	30	37	35	35	00	30	30	30	30	0000755.0000
00000070	30	30	30	00	30	30	30	30	30	30	30	00	30	30	30	30	000.0000000.0000
00000080	30	30	30	30	30	30	30	00	31	34	32	34	37	35	30	33	0000000.14247503
00000090	33	36	36	00	30	31	30	33	34	31	00	20	35	00	00	00	366.010341. 5
029FEC80	00	18	10	06	4C	90	92	6D	31	18	10	06	41	ВЗ	D5	1C	L.'mlA'Õ.
029FEC80 029FEC90	00 86	18 5C	10 37	06 00	4C 00	90 00	92 00	6D 00	31 FB	18 06	10 00		41 01	B3 00	D5 00	1C 00	L.'mlA³Õ.
	-			06 00 00	4C 00 00	90 00 00		•									
029FEC90	86			06 00 00	00	90 00 00 00	00	00	FΒ	06	00	00	01	00	00	00	†\7û
029FEC90 029FECA0	86 6B		37 01	00	00 00	00	00 00	00 20	FB 00	06 00	00 00	00 00	01 00	00	00 00	00 00	†\7û k
029FEC90 029FECA0 029FECB0	86 6B 00	5C 14 10	37 01 00	00 00 00	00 00 01	00 00 00	00 00 00	00 20 00	FB 00 02	06 00 64	00 00 63	00 00	01 00 2F	00	00 00 20	00 00 3E	t\7û kdcu/.>
029FEC90 029FECA0 029FECB0 029FECC0	86 6B 00	5C 14 10	37 01 00	00 00 00	00 00 01	00 00 00 30	00 00 00 30	00 20 00	FB 00 02 35	06 00 64 35	00 00 63 00	00 00 75 7C	01 00 2F 00	00 00 00 4C	00 00 20	00 00 3E 2E	t\7û kdcu/.> 0000755. .L.
029FEC90 029FECA0 029FECB0 029FECC0 029FECD0	86 6B 00 00	5C 14 10 00 00	37 01 00	00 00 00	00 00 01	00 00 00 30	00 00 00 30 30	00 20 00 37 00	FB 00 02 35 31	06 00 64 35 34	00 00 63 00 32	00 00 75 7C	01 2F 00 37	00 00 00 4C	00 00 20 00 30	00 00 3E 2E 33	t\7û kdcu/.> 0000755. .L H0.14247503
029FEC90 029FECA0 029FECB0 029FECC0 029FECD0 029FECE0	86 6B 00 00 1C 33	5C 14 10 00 00 36	37 01 00 05 48 36	00 00 00 30 02 00	00 00 01 30 00 30	00 00 00 30 05 31	00 00 00 30 30 30	00 20 00 37 00	FB 00 02 35 31	06 00 64 35 34	00 00 63 00 32 00	00 00 75 7C	01 2F 00 37	00 00 00 4C 35 20	00 00 20 00 30 3E	00 00 3E 2E 33 5C	t\7ûkdcu/.>0000755. .LH0.14247503 366.010341.5>\

Figure 22. Actual file data in dcu.tar file(top)/data node which is dcu.tar file is stored (bottom).

5.2. File Recovery

5.2.1. VDFS

In the case of VDFS, data were stored in the form of snapshots, enabling the recovery of the deleted files. We could search for the deleted file name in the catalog tree or check the file size and records that did not have an offset address where the data were stored. Figure 23 shows the deleted 8669.png file's record. We confirmed that the file size and data area offset address did not exist. This phenomenon occurs when a file is either modified or deleted. However, because VDFS used a snapshot feature, if a new snapshot was not overwritten after the file was deleted, information on the deleted file existed in the previously saved snapshot. Therefore, it was possible to recover deleted data using the previous snapshot information.

																J File	e size		D	ata o	ffset
1D705C80	00	00	00	00	00	00	00	00	00	00	00	00	28	00	60	01				(. ` .
1D705C90	D2	25	00	00	00	00	00	00	EO	2F	00	00	00	00	00	00	Ò%.		.à/		
1D705CA0				36													8	669.	png		• • •
1D705CB0	00	00	00	00	0C	00	00	00	00	00	00	00	00	00	00	0.0					
1D705CC0	01	00	00	00	00	00	00	00	FF			· ŸŸ	ŸŸŸ	ŸŸŸ							
1D705CD0	A4	81	00	00	89	13	00	00	64	00	00	00	D9	D2	DA	60					
1D705CE0	00	00	00	00	EC	82	70	02	D9	D2	DA	60	00	00	00	00		.ì,p			
1D705CF0	EC	82	70	02	D9	D2	DA	60	00	00	00	00	EC	82	70	02	ì,p	. ÙÒÚ	`	ì	,p.
1D705D00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00					
1D705D10	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00					
1D705D20	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00					

Figure 23. Record of deleted 8669.png file without file size and offset.

Using an editor such as HxD, we could obtain deleted file information from previous snapshots by searching for the file name, or for the file name that we wanted to recover. In the previous saved snapshots, an offset value containing the size and data of the 8669.png file existed, as shown in Figure 24. It was, therefore, possible to recover the deleted file by moving it to the corresponding offset address and copying and storing the data in the corresponding area.

														L		File	size Data offset
1E7AE5A0	00	00	00	00	28	00	60	01	D2	25	00	00	00	00	00	00	(.`.Ò%
1E7AE5B0	EO	2F	00	00	00	00	00	00	02	08	38	36	36	39	2E	70	à/8669.p
1E7AE5C0			00	00				00			00	-				00	ng
1E7AE5D0	6F	39	01	00	00	00	00	00	01	00	00	00	00	00	00	00	09
1E7AE5E0	FF	A4	81	00	00	89	13	00	00	ÿÿÿÿÿÿÿ¤ %							
1E7AE5F0	64	00	00	00	D9	D2	DA	60	00	00	00	00	EC	B5	22	18	dὑο΄΄ìμ".
1E7AE600	D9	D2	DA	60	00	00	00	00	EC	B5	22	18	D9	D2	DA	60	ÜÒÚ`ìμ".ÜÒÚ`
1E7AE610	00	00	00	00	EC	82	70	02	6F	39	01	00	00	00	00	00	ì,p.09
1E7AE620	14	00	00	00	00	00	00	00	DD	B3	00	00	00	00	00	00	ݳ
1E7AE630	14	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
1E7AE640	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	

Figure 24. Records of the 8669.png file stored in snapshot format.

5.2.2. JFFS2

JFFS2 is a journaling file system, so it was also possible to recover deleted files. After deleting the test2.txt file, a partition dump was applied in the present study. Figure 25 shows the inode and dirent node of the test2.txt file. The left side of Figure 25 shows that the version was 0x04 with data prior to the deletion of the file, and the right side shows that the version changed to 0x05 with data after the file was deleted. Therefore, when analyzing the JFFS2 file system, if a deleted file with only a dirent node was found, a search for the inode prior to the file deletion could be conducted, and the deleted file could be recovered.

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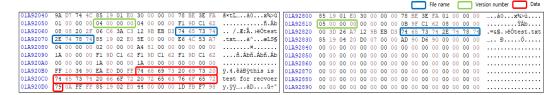


Figure 25. Recovery of deleted files in JFFS2 by verifying file versions.

5.2.3. UBIFS

UBIFS applied compression and encryption when files were stored. As shown in Figure 26, data of dcu.tar file had been encrypted and compressed. Due to this phenomenon, recovery of deleted files in UBIFS was limited.

																	Data
04CDEF60	31	18	10	06	2B	2C	5E	20	CF	5C	37	00	00	00	00	00	1+,^ Ï\7
04CDEF70	92	0B	00	00	01	00	00	00	6B	14	01	00	18	00	00	20	′k
04CDEF80	00	00	00	00	00	00	00	00	00	10	00	00	01	00	00	00	
04CDEF90	00	4F	C4	18	9F	E5	7A	20	82	E2	A7	AB	FF	EB	A0	08	.OÄ.Ÿåz ,â§«ÿë .
04CDEFA0	9F	E5	04	10	A0	El	30	E7	FF	EB	00	10	AO	El	06	20	Ÿå á0çÿë á.
04CDEFB0	A0	El	05	00	A0	El	01	30	A0	E3	F7	AF	00	EB	05	EA	á á.0 ã÷ .ë.ê
04CDEFC0	8D	E2	00	60	A0	El	50	08	9E	E5	0C	40	40	E2	04	00	.â.` áP.žå.@@â
04CDEFD0	57	El	79	01	00	1A	16	4B	8D	E2	4C	40	84	E2	16	2B	WáyK.âL@"â.+
04CDEFE0	8D	E2	04	00	A0	El	74	18	9F	E5	79	20	82	E2	92	AB	.â át.Ÿåy ,â'«
04CDEFF0	FF	EB	4C	CD	0A	18	31	4D	01	E2	6D	0A	2A	AF	0A	4C	ÿëLÍlM.âm.*¯.L
04CDF000	08	92	27	4D	01	23	CD	0A	48	29	4C	01	06	24	18	9F	.''M.#Í.H)L\$.Ÿ
04CDF010	E5	78	20	82	E2	7D	5E	14	F8	07	BD	14	06	2D	4C	01	åx ,â}^.ø.⅓L.
04CDF020	02	03	30	AO	E3	CD	6D	OA	3A	AF	OA	48	08	93	27	4D	0 ãím.: .H."'M

Figure 26. Encrypted data area in dcu.tar file.

Through this experiment, the possibility of extracting and recovering files of three types of file systems was studied. According to the results of the study, the possibility of checking the file name by file system, the possibility of file extraction, and the possibility of recovering deleted files were confirmed, which can be seen in Table 1.

Table 1. The possibility	v of information	available for	each file system.
---------------------------------	------------------	---------------	-------------------

File System	File Name Check	File Extract	File Recovery
VDFS	Able	Able	Able
JFFS2	Able	Able	Able
UBIFS	Able	Disable	Disable

6. Discussion

We conducted experiments on file system analysis, file extraction, and the possibility of recovery of deleted files on IoT devices. Since the experiment used a real device in the experiment, the applied analysis method could also be used in other models.

In this study, file system forensics was conducted for Samsung smart TV. As a result of analyzing Samsung smart TV, it was confirmed that Samsung's own file system, called VDFS, was used. By analyzing metadata, we understood the structure in which the file was stored and the metadata information. Access to the file data area was required to verify file extraction and recoverability of the file system, which was accessible through the previously performed metadata analysis. Studies have shown that VDFS stores data in the form of snapshots, and past snapshots allowed extraction of files and recovery of deleted files.

In addition, file system forensics was conducted on dcu, a power device, as an additional device. As a result of analyzing dcu, it was confirmed that two file systems, JFFS2 and UBFIS, were used. The file systems were also analyzed to confirm metadata and file

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structures and research on the possibility of file extraction and file recovery. JFFS2 uses a journaling system to extract and recover deleted files. However, since the file is compressed and stored, extraction of the file requires decompression, and decompression is possible through a simple python code. UBIFS is also compressed, but since the file is encrypted, the contents of the file cannot be checked even if decompression is performed.

The results of this study can help in the extraction and recovery of files for IoT devices using the same platform. In addition, if only device data can be obtained, the technology used in this study has the advantage of being able to apply forensics quickly anywhere and does not require additional cost or technology. In particular, this paper's research was conducted at the file system stage on VDFS, which was previously conducted only with file recovery research. In addition, we analyzed metadata of some file systems installed in IoTs, which have not been implemented much in the past, and this differentiated our study from studies that focused only on file recovery.

However, some limitations existed in this study. In the case of UBFIS, compression and encryption proceed when the file is saved, but the method for decryption could not be confirmed. Accordingly, there was a limitation in that original data stored in UBIFS could not be collected. It also required analysis of some metadata. Analysis did not proceed because it was difficult to confirm what information some metadata stored, which made it difficult to find the desired file in the directory. For example, additional analysis of metadata, such as the data table, required finding the desired file. Therefore, in a future study, additional metadata will be analyzed to enable faster file navigation. In addition, for the file extraction of UBIFS, research on the decryption of encrypted data areas is needed. In addition, further research will be conducted on platforms and file systems used in small IoT devices.

7. Conclusions

Various IoT platforms are currently in use, and analysis studies applying to IoT platforms are extremely important from a forensic or security maintenance perspective. As various platforms are used, various evidence acquisition approaches are needed in the event of a crime. In this study, we analyzed the file systems VDFS, JFFS2 and UBIFS which are used on the Tizen and Linux platforms. This study presented an approach for conducting a metadata analysis, as well as an approach for data acquisition and recovery when analyzing other file systems for reference. In this study, the file storage structure of a specific file system was analyzed through metadata analysis. It also showed the possibility of extracting stored files by suggesting methods, such as decompression and decryption, according to methods, such as encryption and compression, used when storing files. In addition, when the file was deleted, a method of recovering the deleted file was studied according to the method of storing the file in the file system, such as journaling and snapshots. Finally, the division into file name verification, file extraction, and file recovery stages facilitates checking the applicability of the file system step by step. Through this, more diverse data can be secured from IoT devices and used as evidence in the event of a crime. Moreover, based on the analysis results of this study, future IoT devices with greater security can be developed. Through the metadata analysis, it is expected that the results of this study will contribute to securing evidence in the event of a crime. In particular, since VDFS analysis has not been previously studied, it may be helpful to secure evidence from smart TVs through this study. In addition, it is expected to help secure evidence when investigating IoT devices using file systems with similar structures. Therefore, the metadata analysis on the file system conducted in this study is expected to help improve the overall security technology of the IoT platform, as well as the forensics aspect.

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